



Efficient Resource Management in Wireless Mesh Network

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EU-MESH Project

- FP7 ICT-1-1.1, Call 1 (May 2007)
- Start date, duration: 1 Jan 2008, 30 months
- 9 partners from 7 EU member states:





Partners

1. Foundation for Research and Technology – Hellas (FORTH), GR: **Coordinator**
 2. National Research Council (CNR), IT
 3. Technical University Berlin, DE
 4. SUPSI, CH
 5. Budapest University of Technology & Economics (BME), HU
 6. Proximity Poland, PL
 7. Thales, FR
 8. Hellenic Telecommunications and Telematic Applications Company (Forthnet), GR
 9. Ozone, FR
- } Research/Academic
- } SME, wireless mgmt software
- } Systems integrator/manufacturer
- } Telecom/wireless provider

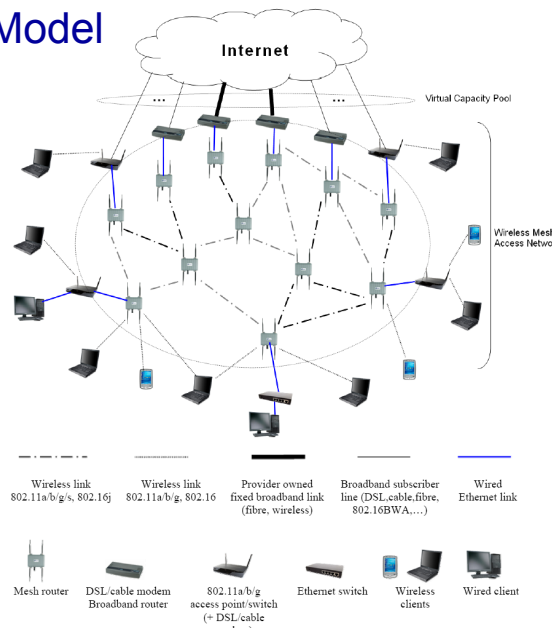
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EU-MESH Network Model

- Wireless mesh network provides access to **virtual capacity pool**
- Virtual capacity pool aggregates capacity from many **subscriber & provider links**
- Access for **stationary and mobile users**
- Support for **QoS, dependability, and security**



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Objectives and key issues

- *Develop, evaluate, and trial a system of software modules for building wireless multi-radio multi-channel mesh networks that*
 - support scalable end-to-end QoS while efficiently utilize wireless spectrum and control interference
 - enable fast deployment and reduce management complexity and cost, while providing dynamic and reliable connectivity
 - support enhanced proactive & reactive security, and seamless mobility based on cross-layer monitoring in single and multi-operator environments
 - seamlessly integrate mesh networks with fixed technologies to provide pervasive ultra-high capacity broadband access to both stationary and mobile users, through a converged infrastructure
- Implement and evaluate system in two existing metropolitan scale deployments (Paris and Heraklion)

Jointly consider QoS, mobility, and security



Key Research & Technical Achievements

- Requirements and Architecture (WP2)
 - Technical requirements and functionalities considering jointly QoS, mobility, and security
 - Initial design of cross-layer architecture
- Mesh Configuration and Link Control (WP3)
 - Hybrid channel access scheme for interference reduction
 - Initial design of auto-configuration & self-healing procedures
 - Initial design of real-time, technology independent cross-layer monitoring system
 - Dynamic approximation of mesh nodes' coverage



Research & Technical Achievements (cont.)

- Resource, QoS, and Mobility Mgmt (WP4)
 - *Channel assignment in metropolitan mesh networks and utility-based joint channel assignment & topology control*
 - *Congestion-aware routing metrics capturing contention, rate diversity, and interference*
 - *Queuing theoretic network model for capacity estimation in multi-homed mesh networks*
- Security (WP5)
 - Framework for **secure link-state routing** protocol
 - Authentication supporting **fast handover in multi-operator** environment and continuous access control enforcement
 - Route segment protection scheme based on IPsec
 - **Anomaly-based intrusion detection** for jamming attacks

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Rest of this talk

- Joint channel assignment & topology control
 - flexibility in having different target objectives
- Channel assignment in metropolitan networks
 - timescales of channel assignment and rate control
- Contention aware routing
 - inter- and intra-flow contention

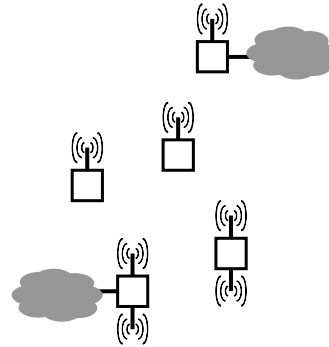
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Joint Channel Assignment & Topology Control

- Model
 - network of mesh nodes with one or more interfaces
 - some mesh nodes connected to fixed network (gateways - GWs)
- Problem: assign channels to achieve some **target objective**
 - multi-objective



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Flexibility for Channel & Topology Control

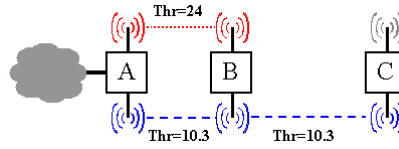
- Node/interface connectivity **not known a priori**
 - jointly perform **channel assignment & topology control**
 - ensures **connectivity to gateway** for all nodes
- Channel assignment based on **different objectives**
 - objectives expressed as **utility functions of throughput**
 - throughput estimation accounts for **rate diversity**
- **Different utility functions** result in **different channel assignments**
 - **aggregate throughput**
 - **fairness** in allocation of throughput among node pairs
 - **redundancy**: multiple links with different interfaces

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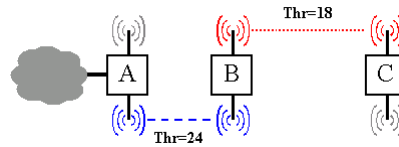
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Comparison of objectives

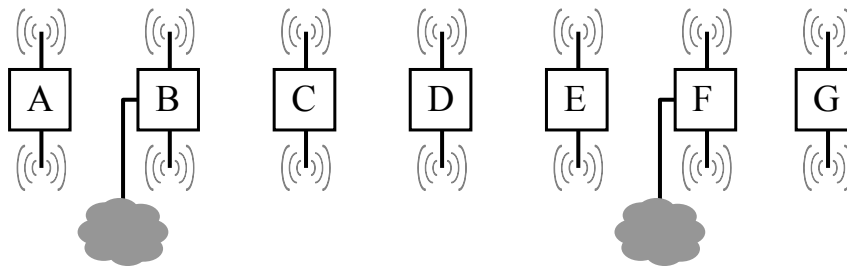
- Throughput objective:
 $\text{Thr}_{A-B}=34.3$, $\text{Thr}_{B-C}=10.3$
 $\text{Thr}_{\text{total}}=44.3$

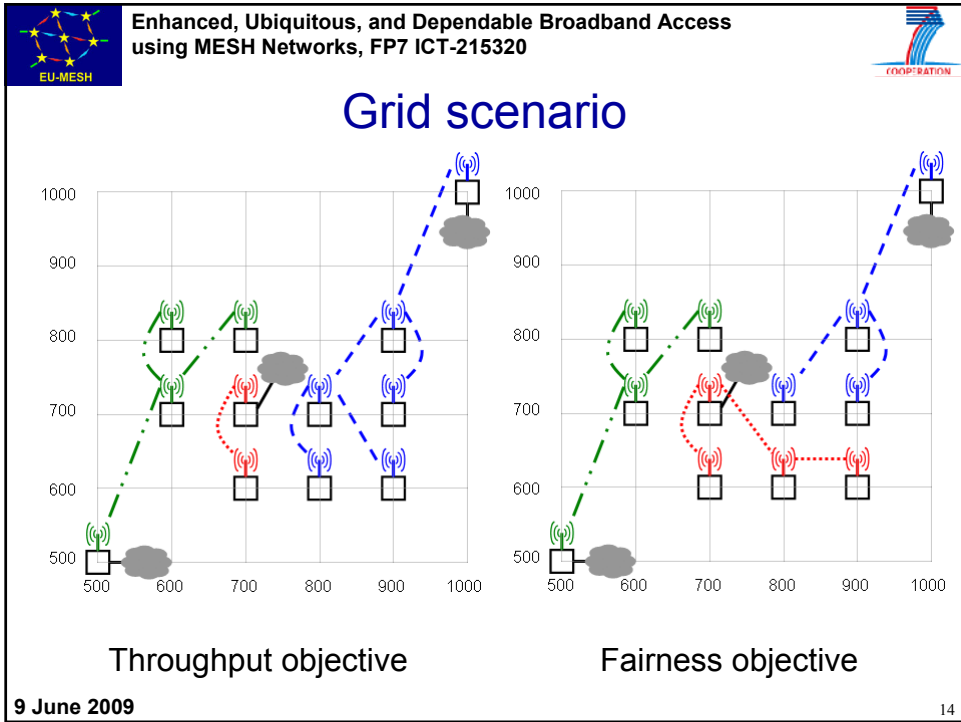
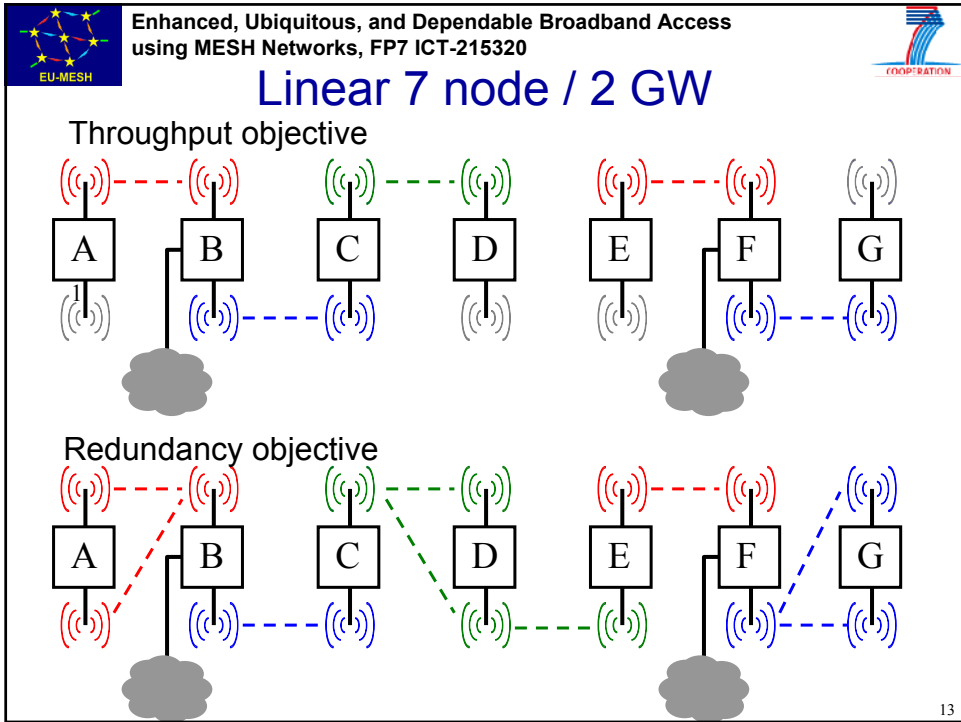


- Fairness objective:
 $\text{Thr}_{A-B}=24$, $\text{Thr}_{B-C}=18$
 $\text{Thr}_{\text{total}}=42$



Linear 7 node / 2 GW

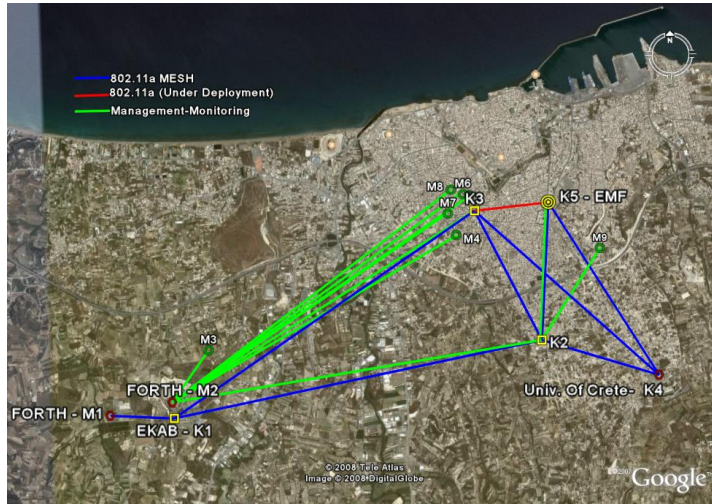






Heraklion Metropolitan MESH Test-bed

- 60 Km² coverage
- 14 nodes, 5 core multi-radio nodes
- 1.6 – 5 Km links
- Antennas: 19/21/26 dBi panel
- Three fixed network connections (FORTH & UoC)
- Independent mgmt / monitoring network



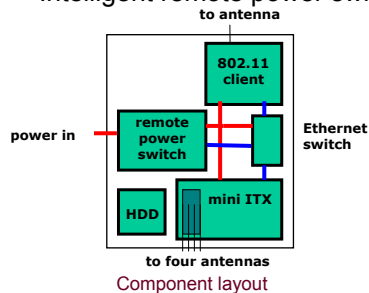
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Multi-radio mesh node

- Mini-ITX board (EPIA SP 13000, 1.3 GHz C3, 512 MB ram)
- Four mini PCI Atheros-based 802.11a/g wireless cards
- Gentoo 2006 i686 Linux (2.6.18), MadWiFi version 0.9.2
- OLSR daemon version 0.4.10
- Independent 802.11a client for management/monitoring
- Intelligent remote power switch



Actual node

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Actual deployment pictures



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Channel assignment algorithm

interference model

- multi-point link conflict graph
- measurement-based interference estimation using test-traffic

link ordering

- distance from fixed network gateway
- increasing SNR
- random

channel selection metric

- one-way SNR
- two-way SNR
- two-way packet delay

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Multi-Point Link Conflict Graph - MPLCG

- Conflict graph for modeling intra-network interference
 - vertices correspond to a link between two nodes
 - in multi-radio nodes with directional antennas, not all interfaces equivalent
- New proposal: Multi-Point Link Conflict Graph
 - vertex represents a multi-point communication link - *set of two or more interfaces belonging to different nodes*
 - *edge between two vertices* indicates that the two corresponding links *interfere with each other*
 - requires *a priori knowledge of the communication graph* which identifies the interfaces that will communicate with each other

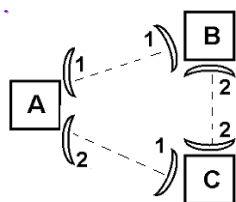
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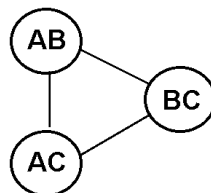


Examples of MPLCG

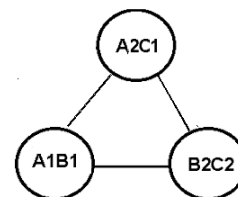
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3-node network

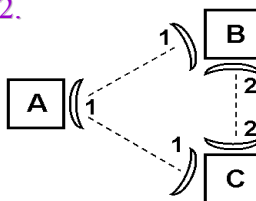


conflict graph

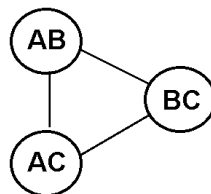


multi-point link conflict graph

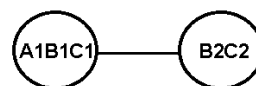
2.



3-node network



conflict graph



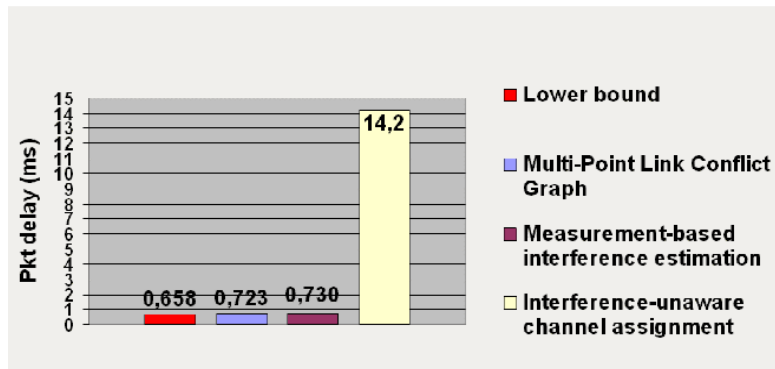
multi-point link conflict graph

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Comparison with lower bound and interference unaware approach



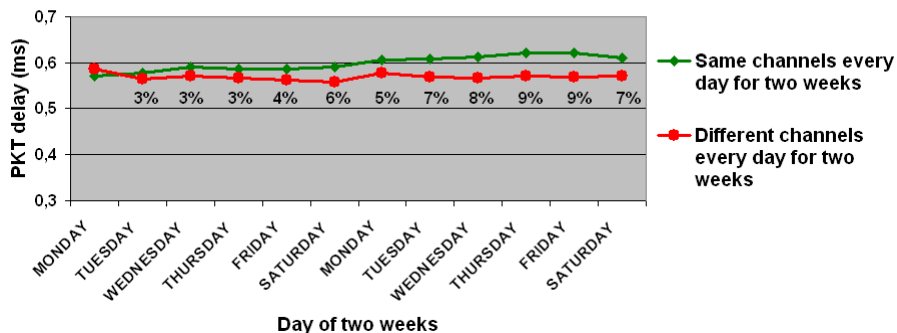
- MPLCG, measurement-based interference estimation: avg pkt delay approximately **11%** higher than lower bound
- interference-unaware scheme: **20 times** higher avg delay

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Same and different channels for two weeks



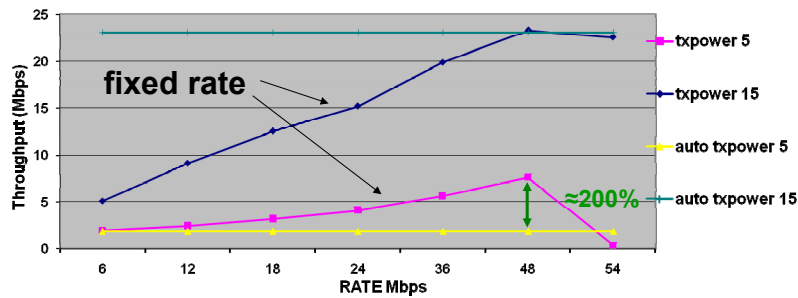
- fixed channel assignment for up to two weeks remains close to channel assignment performed every day
- for time interval of two weeks channel assignment should not be updated

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Rate control in metropolitan links

Link: K2 ↔ K3, distance: 2 Km Avg SNR: 28



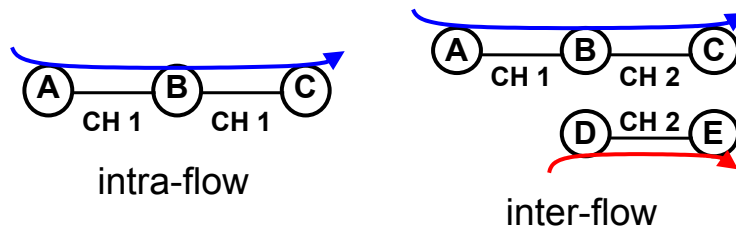
- Low SNR: fixed-rate can achieve significantly higher throughput compared to auto-rate
- High SNR: fixed-rate & auto-rate achieve same performance

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Routing metric in multi-rate WMNs

- Isotonicity: necessary property of routing metric to apply Dijkstra algorithm & hop-by-hop routing
- Contention: intra-flow and inter-flow



- Contention level depends on # of flows and transmission rate

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Approximation for maximum throughput

- Upper bound for throughput at transmitter i

$$x_i = \frac{L_i}{\sum_{j \in N_i} \frac{L_j}{R_j}}$$

- N_i : set of transmitters interfering with i (including i)
- R_j : j 's transmission rate
- L_j : j 's packet size
- Assumption: Saturated transmitters, fair MAC sharing
 - extensions to address these



CATT: Contention-Aware Transmission Time

- CATT metric for link l : packet transmission time estimate based on throughput approximation

$$CATT_l = \sum_{j \in N_l} \frac{L_j}{R_j}$$

- CATT captures influence on link l of all transmitters that interfere with l
- Both inter- and intra-flow interference
- Cost W of path p

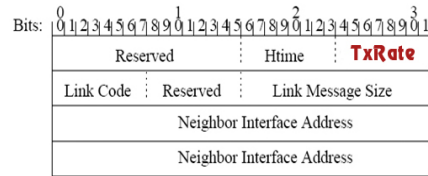
$$W_p = \sum_{l \in p} CATT_l$$

CATT implementation in OLSR

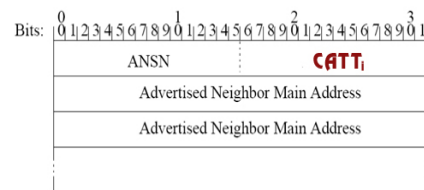
- Assume same pkt size

$$CATT_i = \sum_{j \in N_i} \frac{1}{R_j}$$

- OLSR (Optimized Link State Routing) has two types of messages
- Mesh node obtains TxRate from wireless driver
- Node broadcasts TxRate in Hello messages
- Node computes CATT metric
- Node broadcasts CATT metric in TC messages



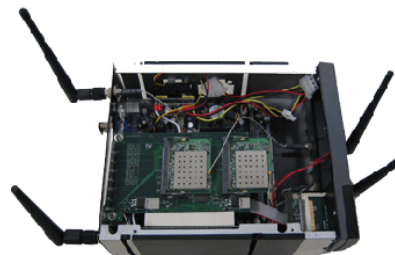
Hello msg header



TC msg header

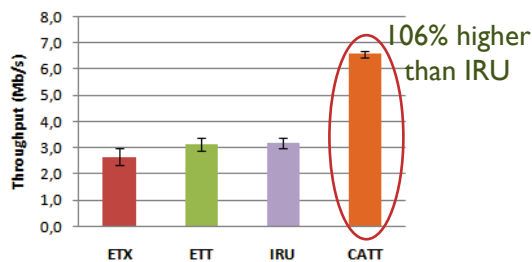
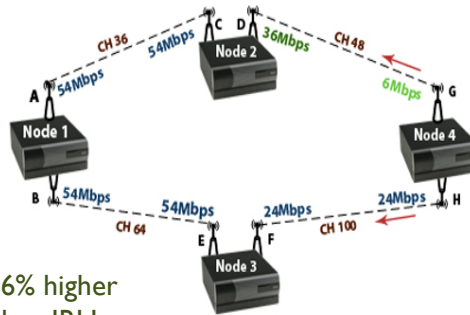
Wireless multi-radio mesh test-bed

- 6 (now 14) four-radio nodes
- Mini-ITX (VIA EPIA Nehemiah M10000G) with 1GHz, 512 MB (DDR400)
- MikroTik Router BOARD 14 4 MiniPCI to PCI adapter
- 4 High Power Super A/G Atheros 802.11a/b/g
- Ubuntu 7.04 Linux kernel 2.6.20-16-generic, MadWiFi 0.9.3.1
- OLSR daemon 0.4.10



Asymmetric link scenario: constant rate traffic

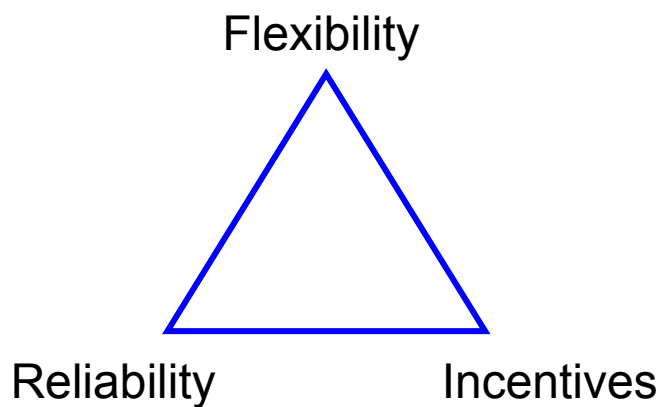
- ETT, ETX & IRU mostly select path 1-2-4: Rate on D is higher (36Mbps) than F (24Mbps)
- CATT captures the interference due to low rate of interface G (6Mbps)



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Three issues key to design and operation
of **wireless access networks**:



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